Preparing seismic codes for GPUs and other many-core architectures

Paulius Micikevicius Developer Technology Engineer, NVIDIA

2010 SEG Post-convention Workshop (W-3) High Performance Implementations of Geophysical Applications Oct 21, 2010

Outline

- High level review of various GPUs
- Common features to all GPUs
- Requirements for maximizing GPU performance
- Requirements as applied to
 - Kirchhoff Migration
 - Reverse Time Migration

Terminology

- We'll define a "brick" for this talk
 - Since all vendors use different names
 - GPUs are built by replicating bricks (10s)
 - Connected to memory via some network
- Brick = minimal building block that contains own:
 - Control unit(s) decodes/issues instructions
 - Registers
 - Pipelines for instruction execution
 - Local cache(s)

NVIDIA GPU Brick (Streaming Multiprocessor)

- Up to 1536 threads
- Instructions are issued per 32 threads (warp)
 - Think 32-way vector of threads
- Source code is for a single thread and is <u>scalar</u>:
 - No vector intrinsics a la SSE
 - HW handles grouping of threads into vectors, vector control flow
- Dual-issue: instructions from different warps
- Shared memory, L1 cache
- Large register file, partitioned among threads

AMD GPU Brick (SIMD Engine)

- Up to ~1500 threads
- Instructions are issued per 64 threads
- VLIW instruction issue:
 - HW designed for 5 "issue" slots (16x5 'cores' per brick)
 - Combine up to 5 instructions from the same thread to maximize performance
- Source code is for a single thread and is <u>scalar</u>:
 - HW handles grouping threads into vectors and control flow
 - Compiler handles VLIW combining
- Shared memory, L1 cache
- Large register file, partitioned among threads

Intel Larrabee/Knights Ferry/Corner Brick (Core)

- Up to 4 threads
- Scalar and vector (512-bit SIMD) units
 - For example: 16-fp32 vector SIMD
- Dual issue: scalar-vector, from the same thread
- Source code is for a single thread and is <u>vector</u>:
 - Intrinsics for SIMD operations (a la SSE)
- L1 and L2 caches
 - Intrinsics for pre-fetching and prioritization of cache lines
 - No user-managed shared memory
- Small register file (relies on caches)

HW Commonalities

- Built by replicating 10s of "bricks"
 - In-order instruction issue
- High GPU memory bandwidth (150+ GB/s)

• "Bricks" are vector processors

- Different execution paths within vectors are supported but degrade performance
- Different execution paths in different vectors have no impact on performance

Vectors access memory in cache-lines

- Consecutive threads (vector elements) should access a contiguous memory region
- Scattered access is supported, but will fetch multiple lines, increasing bandwidth-pressure

Requirements for Maximum Performance

• Have sufficient parallelism

- At least a few 1,000 of threads per function

• Coherent memory access

- By threads in the same "thread-vector"

Coherent execution

- By threads in the same "thread-vector"

Amount of Parallelism

- GPUs issue instructions in order
 - Issue stalls when instruction arguments are not ready
- GPUs switch between threads to hide latency
 - Context switch is free: thread state is partitioned (large register file)
- Conclusion: need enough threads to hide math latency and to saturate the memory bus
 - Independent instructions within a thread also help
- Very rough rule of thumb:
 - Need ~512 threads per "brick"
 - So, at least a few 1,000 threads per GPU

Memory Access

- Addresses from a warp ("thread-vector") are converted into line requests
 - NVIDIA line sizes: 32B and 128B
 - Goal is to maximally utilize the bytes in these lines



Memory Access



Memory Access



Coherent Execution



Execution within warps is coherent



instructions / time

Execution diverges within a warp



Requirements for Maximum Performance

- Have sufficient parallelism
- Coherent memory access
- Coherent execution

Kirchhoff Migration

One (very) simplified way to look at it:

for each input trace (src-rcv pair) do for each output point do compute travel-times get input trace value(s) based on travel times update the output point

Kirchhoff Migration

• Amount of parallelism

- Plenty of output points
- For example: 1 or several points per thread

Coherent memory access

- Successive threads should process adjacent output points
 - Fastest varying dimension for thread IDs = fastest varying dimension for data
 - Coherent writes
 - Reads will be scattered, but within close proximity

Coherent execution

- Usually all threads will execute the same code path

Reverse Time Migration

- The main component is FD computation
 - Here I assume we work in time domain
- Derivatives are computed with stencils



Reverse Time Migration

- The main component is FD computation
- Derivatives are computed with stencils

for each FD step do for each output point do compute wave-field derivatives combine derivatives with property-fields update the output point

Reverse Time Migration

• Amount of parallelism

- Plenty of output points (100s of millions)
- For example: 1 or several points per thread

Coherent memory access

- Successive threads in a warp should process adjacent points
 - Fastest varying dimension for thread IDs = fastest varying dimension for data
 - 2D blocks of threads -> exploit wavefield locality from SMEM
 - Coherent writes and reads
 - Regular grid -> regular reads as well

Coherent execution

- All threads execute the same code path
 - Execution is not data dependent

Summary of Coding Patterns for GPUs

- Seismic problems have plenty of parallelism
 - Output points, traces/shots
 - Need to make sure that the framework exposes that when calling functions parallelizable on accelerators

• Memory accesses can be made coherent enough

- Access patterns are often coherent, or scattered in reasonably close proximity (for L1 or tex cache on GPU)
- Generally, make sure that thread-ID varying-order matches data varying-order
 - If arranging threads into 2D blocks, fastest varying dimension should be accessing in multiples of cache lines

• Strive for coherent execution

Preferably group together threads (hence data) that follows the same path

Comparing GPU and CPU requirements

Amount of parallelism

- CPUs need less currently, but requirement is growing:
 - Increasing core count
 - Increasing vector width

• Coherent memory access

- Already needed to maximize SSE throughput
- Scattered accesses also underutilize bus bandwidth

Coherent execution

Already needed within SSE vectors

References

- NVIDIA:
 - <u>http://www.nvidia.com/object/fermi_architecture.html</u>
 - Fermi Compute Architecture Whitepaper (<u>http://www.nvidia.com/content/PDF/fermi_white_papers/NVIDIA_Fe_rmi_Compute_Architecture_Whitepaper.pdf</u>)
- AMD:
 - Unleashing the Power of Parallel Compujte with Commodity ATI Radeon 5800 GPU, Siggraph Asia 2009 (<u>http://sa09.idav.ucdavis.edu/docs/SA09_AMD_IHV.pdf</u>)
- Intel:
 - L. Seiler et al. Larrabee: A Many Core X86 Architecture for Visual Computing. Siggraph 2008 (<u>http://software.intel.com/file/18198/</u>)
 - Tom Forsyth. The Challenge of Larrabee as GPU. Colloquium at Stanford, 2010

(http://www.stanford.edu/class/ee380/Abstracts/100106.html)

References

- "Seismic Imaging" by S. Morton
 - <u>http://gpgpu.org/wp/wp-</u> content/uploads/2009/11/SC09 Seismic Hess.pdf
- "Accelerating Kirchhoff Migration by CPU and GPU Cooperation", J. Panetta *et al*
 - <u>http://www.cos.ufrj.br/~monnerat/SBAC_2009.html</u>
- "Implementing 3D Finite Difference Codes on the GPU" by P. Micikevicius
 - Slides: <u>http://www.nvidia.com/content/GTC/documents/1006_GTC09.</u> <u>pdf</u>
 - video (slides+audio): <u>http://developer.download.nvidia.com/compute/cuda/docs/GT</u> <u>C09Materials.htm</u> then look for talk 1006

BACKUP

Current Differences among GPUs

• Hardware:

- Vector width: 16, 32, 64
- VLIW (AMD) vs dual-issue (NVIDIA, Intel)
- Dual-issue: same thread (Intel) vs different threads (NVIDIA)
- Large register file, small cache (NVIDIA, AMD) vs small register file, larger caches (Intel)

• Programming model

- Intel: vector source code (SIMD intrinsics)
- AMD, NVIDIA: scalar source code
 - hw aggregates threads into vectors and resolves control flow divergence

Simple 2D Stencil Code in CUDA (OpenCL equivalent in comments)

__global__ void stencil_2d(float *output, float *input, const int dimx, const int dimy, const int row_size)

int ix = blockIdx.x*blockDim.x + threadIdx.x; // ix = get_global_id(0); int iy = blockIdx.y*blockDim.y + threadIdx.y; // iy = get_global_id(1); int idx = iy * row_size + ix;

```
output[idx] = -4 * input[idx]
+ input[idx+1] + input[idx-1]
+ input[idx + row_size] + input[idx - row_size];
```

{

}